

# Lattice reduction and continued fractions

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# Continued fractions

We consider a positive real number  $\alpha$ .

One looks for sequences of rational numbers  $(p_n/q_n)_n$  that satisfies

$$\lim p_n/q_n = \alpha$$

Continued fractions allow to do it with **exponential speed**

$$|\alpha - p_n/q_n| \leq \frac{1}{q_n^2}$$

## Continued fractions

We represent real numbers in  $(0, 1)$  as

$$\alpha = \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \frac{1}{a_4 + \dots}}}}$$

with the **partial quotients** (digits)  $a_i$  being **positive integers**

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Rational approximations are then given by

$$p_n/q_n = \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{\dots + \frac{1}{a_n}}}} \quad \left| \alpha - p_n/q_n \right| \leq \frac{1}{q_n^2}$$

# Dirichlet's bound and exponential convergence

## Dirichlet's theorem

Given real numbers  $(\alpha_1, \dots, \alpha_d)$ , for any positive integer  $N$ , there exist integers  $p_1, \dots, p_d, q$  with

$$1 \leq q \leq N$$

such that

$$\left| \frac{p_i}{q} - \alpha_i \right| < \frac{1}{q N^{1/d}} \quad i = 1, 2, \dots, d$$

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such that

$$\left| \frac{p_i}{q} - \alpha_i \right| < \frac{1}{q N^{1/d}} \leq \frac{1}{q^{1+\frac{1}{d}}} \quad i = 1, 2, \dots, d$$

Dirichlet's bound  $1 + 1/d$

# Euclid algorithm

We start with two nonnegative integers  $u_0$  and  $u_1$

$$u_0 = u_1 \left[ \frac{u_0}{u_1} \right] + u_2$$

$$u_1 = u_2 \left[ \frac{u_1}{u_2} \right] + u_3$$

$\vdots$

$$u_{m-1} = u_m \left[ \frac{u_{m-1}}{u_m} \right] + u_{m+1}$$

$$u_{m+1} = \gcd(u_0, u_1)$$

$$u_{m+2} = 0$$

One **subtracts** the smallest number to the largest as much as we can

# Euclid algorithm and continued fractions

We start with two coprime integers  $u_0$  and  $u_1$

$$u_0 = u_1 a_1 + u_2$$

$\vdots$

$$u_{m-1} = u_m a_m + u_{m+1}$$

$$u_m = u_{m+1} a_{m+1} + 0$$

$$u_{m+1} = 1 = \gcd(u_0, u_1)$$



# Euclid algorithm and continued fractions

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$$\frac{u_1}{u_0} = \frac{1}{a_1 + \frac{u_2}{u_1}}$$

$$u_1/u_0 = \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{\ddots + \frac{1}{a_m + \frac{1}{a_{m+1}}}}}}$$

## Matricial description

We start with two real numbers  $(x_0, x_1)$  in  $(0, 1)^2$  with  $x_0 > x_1$

We divide the largest entry by the smallest and we continue

$$x_0 = \lfloor x_0/x_1 \rfloor x_1 + x_2 \qquad a_1 := \lfloor x_0/x_1 \rfloor$$

$$\begin{pmatrix} x_0 \\ x_1 \end{pmatrix} = \begin{pmatrix} a_1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} a_1 & 1 \\ 1 & 0 \end{pmatrix} \cdots \begin{pmatrix} a_n & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} x_n \\ x_{n+1} \end{pmatrix}$$

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We normalize  $\alpha := x_1/x_0$  and we set

$$M_n := \begin{pmatrix} a_n & 1 \\ 1 & 0 \end{pmatrix} \rightsquigarrow \begin{pmatrix} 1 \\ \alpha \end{pmatrix} \in \bigcap_n M_1 \cdots M_n \mathbb{R}_+^2$$

$$M_1 \cdots M_n = \begin{pmatrix} q_n & q_{n-1} \\ p_n & p_{n-1} \end{pmatrix} \rightsquigarrow \text{a sequence of lattice basis for } \mathbb{Z}^2$$

# Multidimensional continued fractions

If we start with two parameters  $(\alpha, \beta)$ , one looks for two sequences of rational numbers  $(p_n/q_n)$  et  $(r_n/q_n)$  with the **same denominator** that satisfy

$$\lim p_n/q_n = \alpha \quad \lim r_n/q_n = \beta$$

Expected speed 3/2

$$|\alpha - p_n/q_n| \leq 1/q_n^{3/2} \quad |\beta - r_n/q_n| \leq 1/q_n^{3/2}$$

# Canonicity of continued fractions

- **Euclid's algorithm** Starting with two numbers, one subtracts the smallest to the largest
- **Unimodularity**

$$\det \begin{pmatrix} q_{n+1} & q_n \\ p_{n+1} & p_n \end{pmatrix} = \pm 1$$

- **Best approximation property**

**Theorem** A rational number  $p/q$  is a **best approximation** of the real number  $\alpha$  if every  $p'/q'$  with  $1 \leq q' \leq q$ ,  $p/q \neq p'/q'$  satisfies

$$|q\alpha - p| < |q'\alpha - p'|$$

Every best approximation of  $\alpha$  is a **convergent**

## From $SL(2, \mathbb{N})$ to $SL(3, \mathbb{N})$

Rem  $SL(2, \mathbb{N})$  is a **finitely generated free** monoid. It is generated by

$$\begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \text{ and } \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

- $SL(2, \mathbb{N})$  is a **free and finitely generated** monoid
- $SL(3, \mathbb{N})$  is not free
- $SL(3, \mathbb{N})$  is not finitely generated. Consider the family of matrices

$$\begin{pmatrix} 1 & 0 & n \\ 1 & n-1 & 0 \\ 1 & 1 & n-1 \end{pmatrix}$$

These matrices are **undecomposable** for  $n \geq 3$  [Rivat]

# Multidimensional continued fractions

There is no **canonical generalization** of continued fractions to higher dimensions

Several approaches are possible

- **Best simultaneous approximations**  
Every  $q'$  with  $1 \leq q' < q$  satisfies  $|||q(\alpha, \beta)||| < |||q'(\alpha, \beta)|||$   
But we lose unimodularity, and the sequence of best approximations depends on the chosen norm [Lagarias]
- **Klein polyhedra and sails** [Arnold]
- **Unimodular** multidimensional Euclid's algorithms
  - sequences of **nested cones** approximating a direction  
Jacobi-Perron algorithm, Brun algorithm [Brentjes, Schweiger]
  - lattice reduction (LLL)  
[Lagarias], [Ferguson-Forcade], [Just],  
[Grabner-Lagarias][Bosma-Smeets][Beukers]

# What is expected?

We are given  $(\alpha_1, \dots, \alpha_d)$  which produces a sequence of basis of  $\mathbb{Z}^{d+1}$  and/or a sequence of approximations

**Arithmetics** A two-dimensional continued fraction algorithm is expected to

- detect integer relations for  $(1, \alpha_1, \dots, \alpha_d)$
- give algebraic characterizations of periodic expansions
- converge sufficiently fast
- provide good rational approximations

**Good** means “with respect to **Dirichlet's theorem**”: there exist infinitely many  $(p_i/q)_{1 \leq i \leq d}$  such that

$$\max_i |\alpha_i - p_i/q| \leq \frac{1}{q^{1+1/d}}$$



We also want...

- to understand generic behaviour
- to be able to control the number of executions if the parameters are rational etc.

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- to understand generic behaviour

### Continued fractions

$$\lim \frac{\log q_n}{n} = \frac{\pi^2}{12 \log 2} = 1.18\dots \quad \text{for a.e. } \alpha$$

$$\lim \frac{1}{n} \{k \leq n; a_k = a\} = \frac{1}{\log 2} \log \frac{(k+1)^2}{k(k+2)} \quad \text{for a.e. } \alpha$$

- to be able to control the number of executions if the parameters are rational etc.

### Continued fractions

$\ell(u, v)$ : number of steps in Euclid algorithm  $0 < v < u$

For  $0 < v < u \leq N$  and  $\gcd(u, v) = 1$

$$\mathbb{E}_N(\ell) \sim \frac{12 \log 2}{\pi^2} \cdot \log N \quad \text{average case}$$

# Multidimensional Euclid's algorithms: a zoo of algorithms

- **Jacobi-Perron** [Jacobi'1868–Perron'1907]: we subtract the first one to the two other ones with  $0 \leq x_1, x_2 \leq x_3$

$$(x_1, x_2, x_3) \mapsto (x_2 - \lfloor \frac{x_2}{x_1} \rfloor x_1, x_3 - \lfloor \frac{x_3}{x_1} \rfloor x_1, x_1)$$

- **Brun** [Brun'1919]: we subtract the second largest and we reorder with  $x_1 \leq x_2 \leq x_3$

$$(x_1, x_2, x_3) \mapsto (x_1, x_2, x_3 - x_2)$$

- **Poincaré**: we subtract the previous one and we reorder with  $x_1 \leq x_2 \leq x_3$

$$(x_1, x_2, x_3) \mapsto (x_1, x_2 - x_1, x_3 - x_2)$$

- **Selmer**: we subtract the smallest to the largest and we reorder with  $x_1 \leq x_2 \leq x_3$

$$(x_1, x_2, x_3) \mapsto (x_1, x_2, x_3 - x_1)$$

- **Fully subtractive**: we subtract the smallest one to all the largest ones and we reorder with  $x_1 \leq x_2 \leq x_3$

# Poincaré algorithm [Nogueira'95]

$$(x_1, x_2, x_3) \mapsto (x_1, x_2 - x_1, x_3 - x_2), \quad x_1 \leq x_2 \leq x_3$$

$$1/\varphi^2 + 1/\varphi = 1$$

$1/\varphi^2$	$1/\varphi$	100
$1/\varphi^3$	$1/\varphi^2$	$100 - 1/\varphi$
$1/\varphi^4$	$1/\varphi^3$	$100 - 1/\varphi - 1/\varphi^2$
$\dots$	$\dots$	$\dots$
$1/\varphi^{k+1}$	$1/\varphi^k$	$100 - \sum_{i < k} 1/\varphi^i$

# Jacobi-Perron algorithm

## Continued fractions

$$\begin{pmatrix} 1 \\ \alpha \end{pmatrix} = \lambda_n \begin{pmatrix} a_1 & 1 \\ 1 & 0 \end{pmatrix} \cdots \begin{pmatrix} a_n & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ \alpha_n \end{pmatrix}$$

## Jacobi-Perron algorithm

$$\begin{pmatrix} 1 \\ \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} 0 & 1 & k \\ 1 & 0 & \ell \\ 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} \beta - \ell\alpha \\ 1 - k\alpha \\ \alpha \end{pmatrix} \text{ with } \ell = \lfloor \beta/\alpha \rfloor, k = \lfloor 1/\alpha \rfloor$$

$$\begin{pmatrix} 1 \\ \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} 0 & 1 & k_1 \\ 1 & 0 & \ell_1 \\ 0 & 0 & 1 \end{pmatrix} \cdots \begin{pmatrix} 0 & 1 & k_n \\ 1 & 0 & \ell_n \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ \alpha_n \\ \beta_n \end{pmatrix}$$

$$\begin{pmatrix} 1 \\ \alpha \\ \beta \end{pmatrix} = \lambda_n \begin{pmatrix} q_n & q'_n & q''_n \\ p_n & p'_n & p''_n \\ r_n & r'_n & r''_n \end{pmatrix} \begin{pmatrix} 1 \\ \alpha_n \\ \beta_n \end{pmatrix}$$

# Unimodular multidimensional continued fractions

A **unimodular**  $d$ -dimensional **continued fraction** map over  $[0, 1]^d$  is a map  $T : [0, 1]^d \rightarrow [0, 1]^d$  such that for any  $\alpha \in [0, 1]^d$ , there is a matrix  $M(\alpha)$  in  $GL(d, \mathbb{Z})$  satisfying

$$\alpha = M(\alpha)T(\alpha)$$

The associated continued fraction algorithm consists in iteratively applying the map  $T$  on a vector  $\alpha \in [0, 1]^d$ . This yields the following sequence of matrices, called the **continued fraction expansion** of  $\alpha$

$$(M(T^n(\alpha)))_{n \in \mathbb{N}}.$$

Set  $M_n := M(T^n(\alpha))$

$$\alpha = M_1 \cdots M_n T^n(\alpha)$$

If the matrices have nonnegative entries, the algorithm is said to be **nonnegative** (Perron–Frobenius theory)

# About nonnegative matrices

Theorem of Perron–Frobenius type [Furstenberg]

One considers an infinite product of matrices

$$E_1 \cdots E_k \cdots$$

with entries in  $\mathbb{N}$ . One assumes that there exists a matrix  $B$  with **strictly positive entries** s.t. there exist

$i_1 < j_1 < \cdots < i_k < j_k$  s.t.

$$B = E_{i_1} \cdots E_{j_1}, \dots, B = E_{i_k} \cdots E_{j_k}, \dots$$

Then, the intersection of the cones

$$\bigcap_k E_1 \cdots E_k(\mathbb{R}_+^n)$$

is unidimensional.

Convergence speed? Type of convergence? Weak? strong?

# Convergence

$$\alpha = M_1 \cdots M_n T^n(\alpha) \text{ with } M_1 \cdots M_n = \begin{pmatrix} q_1^{(n)} & \cdots & q_{d+1}^{(n)} \\ p_{1,1}^{(n)} & \cdots & p_{1,d+1}^{(n)} \\ & \cdots & \\ p_{d,1}^{(n)} & \cdots & p_{d,d+1}^{(n)} \end{pmatrix}$$

One considers simultaneous approximations  $\left( \frac{p_{1,j}^{(n)}}{q_j^{(n)}}, \dots, \frac{p_{d,j}^{(n)}}{q_j^{(n)}} \right)$

Weak convergence Convergence in angle

$$\lim_{n \rightarrow +\infty} \left( \frac{p_{1,j}^{(n)}}{q_j^{(n)}}, \dots, \frac{p_{d,j}^{(n)}}{q_j^{(n)}} \right) = (\alpha_1, \dots, \alpha_d)$$

Strong convergence Convergence in distance

$$\lim_{n \rightarrow +\infty} |q_j^{(n)} \alpha_i - p_{i,j}^{(n)}| = 0 \text{ for all } i, j$$



# Convergence of Jacobi-Perron algorithm

**Theorem** There exists  $\delta > 0$  s.t. for almost every  $(\alpha, \beta)$

$$|\alpha - p_n/q_n| < \frac{1}{q_n^{1+\delta}}, \quad |\beta - r_n/q_n| < \frac{1}{q_n^{1+\delta}}$$

where  $p_n, q_n, r_n$  are produced by either by **Brun/Jacobi-Perron** algorithm

Brun [Ito-Fujita-Keane-Ohtsuki'96]

Jacobi-Perron [Broise-Guivarc'h'99]

# Lyapunov exponents

$$A_n(x) = \begin{pmatrix} q_n & q_{n-1} \\ p_n & p_{n-1} \end{pmatrix}$$

Theorem For a.e.  $x$ ,

$$\lim \frac{1}{n} \log q_n = \frac{\pi^2}{12 \log 2} = 1.18 \dots = \lambda_1$$

$\lambda_1$  is the **first Lyapunov exponent**

**First Lyapunov exponent** = "log largest eigenvalue"  $\rightsquigarrow$  size of the matrices/convergents  $A_n(x) \sim q_n(x) \sim e^{\lambda_1 n}$

Number of steps in Euclid's algorithm = size / log eigenvalue

$$\log N / \lambda_1$$

**Second Lyapunov exponent** = "log of the second eigenvalue"  
 $\rightsquigarrow$  measures the distance between column vectors

# Exponentiation based on Brun algorithm

An SPA resistant exponentiation based on Brun's gcd algorithm and addition chains [B.-Plantard]

One performs an exponentiation  $g^e$  for a given  $e$  for a generic group

Cut  $e$  into  $d$  blocks

$$e = \sum_{i=0}^{d-1} e_i 2^{ik/d}$$

For Brun algorithm, as the dimension  $d$  increases, the probability that partial quotients equal to 1 tends to 1 (one performs subtractions and not divisions) [B.-Lhote-Vallée]

↪ Apply Brun algorithm and addition chains

## Higher-dimensional case

Numerical experiments indicate that classical multidimensional continued fraction algorithms seem to cease to be **strongly convergent** for high dimensions. The only exception seems to be the Arnoux-Rauzy algorithm which, however, is defined only on a set of measure zero [[B.-Steiner-Thuswaldner](#)]

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$d$	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$	$d$	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$
2	-0.11216	1.3683	7	-0.01210	1.0493
3	-0.07189	1.2203	8	-0.00647	1.0283
4	-0.04651	1.1504	9	-0.00218	1.0102
5	-0.03051	1.1065	10	+0.00115	0.9943
6	-0.01974	1.0746	11	+0.00381	0.9799

**Table:** Heuristically estimated values for the second Lyapunov exponent and the uniform approximation exponent of the Brun Algorithm

# Multidimensional continued fraction algorithms

- Allowed operations on numbers

$$+, -, /, \times, [], \geq$$

- Allowed operations on matrices: elementary basis transformations
  - interchanging two vectors  $\rightsquigarrow$  permutation matrices
  - adding an integer multiple of one basis vector to another basis vector  $\rightsquigarrow$  transvection matrices

Ex. LLL algorithm Size reduction steps and exchange steps  
Decisions are taken with respect to quadratic norms

LLL approach

# Lattice reduction algorithms

Lattice reduction is based on the following elementary basis transformations on the vectors of the basis  $(b_1, \dots, b_{d+1})$

- **size reduction** the vector  $b_i$  is replaced by  $b_i - \lambda b_j$ ,  $1 \leq j < i$
- **swaps** one exchanges  $b_i$  and  $b_{i+1}$

These operations are decided with respect to the Gram-Schmidt orthogonalization of the basis  $b$

$$b_i^* = b_i - \sum_{j=1}^{i-1} \mu_{i,j} b_j^* \quad \mu_{i,j} = \frac{\langle b_i, b_j^* \rangle}{\langle b_j^*, b_j^* \rangle}$$

- **Size reduction**  $|\mu_{i,j}| \leq 1/2$  for  $i > j$
- **Lovász condition**  $(\delta - \mu_{i+1,i}^2) \|b_i^*\|^2 \leq \|b_{i+1}^*\|^2$



## From lattice reduction to continued fractions

In a letter to Jacobi in 1850, Hermite explained the following idea

Consider

$$\begin{pmatrix} 1 & 0 & \cdots & 0 & -\alpha_1 \\ 0 & 1 & \cdots & 0 & -\alpha_2 \\ \cdots & \cdots & \cdots & \cdots & \cdots \\ 0 & 0 & \cdots & 1 & -\alpha_d \\ 0 & \cdots & \cdots & 0 & t \end{pmatrix}$$

Let  $t > 0$ . We take the corresponding lattice of  $\mathbb{R}^{d+1}$

$$\mathbb{Z}\mathbf{e}_1 + \cdots + \mathbb{Z}\mathbf{e}_d + \mathbb{Z}(t\mathbf{e}_{d+1} - (\alpha_1\mathbf{e}_1 + \cdots + \alpha_d\mathbf{e}_d))$$

A vector of the lattice is of the form

$$\sum_{i=1}^d (p_i - q_t \alpha_i) \mathbf{e}_i + q_t \mathbf{e}_{d+1}$$

Take a **short vector** in  $\Lambda_t$

# How does LLL produce good approximations?

Let

$$M_t := \begin{pmatrix} 1 & 0 & \cdots & 0 & -\alpha_1 \\ 0 & 1 & \cdots & 0 & -\alpha_2 \\ \cdots & \cdots & \cdots & \cdots & \cdots \\ 0 & 0 & \cdots & 1 & -\alpha_d \\ 0 & \cdots & \cdots & 0 & t \end{pmatrix}$$

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- We take  $t$  small
- One has  $\det(M_t) = t$

**Rem:** One changes the lattice at each step instead of changing the bases of a fixed lattice

The parameter  $t$  only occurs in the last line

# How does LLL produce good approximations?

Let

$$M_t := \begin{pmatrix} 1 & 0 & \cdots & 0 & -\alpha_1 \\ 0 & 1 & \cdots & 0 & -\alpha_2 \\ \cdots & \cdots & \cdots & \cdots & \cdots \\ 0 & 0 & \cdots & 1 & -\alpha_d \\ 0 & \cdots & \cdots & 0 & t \end{pmatrix}$$

LLL produces in **polynomial time** a vector  $b_1$  such that

$$\|b_1\| \leq 2^{d/4} \det(M_t)^{1/d+1} = 2^{d/4} t^{1/d+1}$$

One has

$$b_1 = (p_1 - q\alpha_1)e_1 + \cdots + (p_d - q\alpha_d)e_d + qte_{d+1}$$

$$\forall i, \quad |p_i - \alpha_i q| \leq 2^{d/4} t^{1/d+1} \quad \text{and} \quad qt \leq 2^{d/4} t^{1/d+1}$$

$$\rightsquigarrow \forall i, \quad |p_i - \alpha_i q| \leq 2^{(d+1)/4} 1/q^{1/d}$$

# Approximations and lattices

Let

$$M_t := \begin{pmatrix} 1 & 0 & \cdots & 0 & -\alpha_1 \\ 0 & 1 & \cdots & 0 & -\alpha_2 \\ \cdots & \cdots & \cdots & \cdots & \cdots \\ 0 & 0 & \cdots & 1 & -\alpha_d \\ 0 & \cdots & \cdots & 0 & t \end{pmatrix}$$

[Lagarias'93]

$$\forall \alpha = (\alpha_1, \dots, \alpha_d), \forall Q, \exists q, 1 \leq q \leq Q, |||q\alpha||| < \sqrt{d+1}Q^{-1/d}$$

[Lagarias'85, '94, Grabiner-Lagarias'2001]

[Lagarias'94] Let  $t$  tend to 0 and consider Minkowski reduction. The conditions are linear in  $\sqrt{t}$  but when  $n = 7$ , the number of inequalities is about 90,000 for Minkowski reduction.

[Bosma-Smeets'2013] Decrease the value of  $t$  by dividing it by a fixed constant.

[Beukers'2014]

Proves the linearity in  $\sqrt{t}$  of the conditions in LLL.

The values of  $t > 0$  for which  $M_t$  is LLL-reduced form an interval  $[t_0, t_1]$ .

If  $\alpha \notin \mathbb{Q}^d$ , the sequence of critical points is an infinite sequence descending to 0.

# Toward continued fractions

One has  $t \downarrow 0$

- How to change  $t$ ?
- How much does one have to recompute when one changes  $t$ ?
- How to choose stopping times for  $t$ ?
- Can we get nonnegative matrices?
- What are the rules that provide exponential convergence?
- Can we evaluate the growth of the size of the matrices  $M_1 \cdots M_n$ ?